Outline

•Loops

A Loop with a variable array index

Loop: g = g + A[i];

i = i + j;

If (i != h) go to Loop

Variables g, h, i, and j associated with $\$s1 \rightarrow \$s4$ Array base is in \$s5. What is the MIPS assembly code?

Loop:	add \$t1, \$s3, \$s3	# \$t1 ← 2 *i
	add \$t1, \$t1, \$t1	# \$t1 ← 4 *i
	add \$t1, \$t1, \$s5	# \$t1 ← address of A[i]
	lw \$t0, 0(\$t1)	# \$t0 ← A[i]
	add \$s1, \$s1, \$t0	# g ← g + A[i]
	add \$s3, \$s3, \$s4	$\#\ i=i+\ j$
	bne \$s3, \$s2, Loop	# go to Loop if i != h

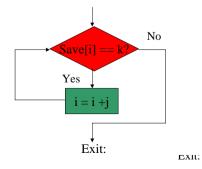
2

A While Loop_{1/2}

while (save[i] == k)

i = i + j;

i, j, and k associated with $s_3 \rightarrow s_5$. Base of array save is in s_6 . What is MIPS Assembly code?



A less than test

if (a < b) go to Less

a, b associated with \$s0, and \$s1, respectively. What is the MIPS assembly code?

•To test for less than we introduce a new instruction

slt \$t0, \$s0, \$s1

slt is set on less than

Register \$t0 is set to 1 if the value in register \$s0 is less than the value in register \$s1, otherwise, register \$t0 is set to 0

•Register \$zero always contains zero

slt \$t0, \$s0, \$s1

•Comparing \$t0 to \$zero gives us the effect of branching if a less than b

•Combining slt with bne implements branch on less than

$t_0 \leftarrow 1$ if a < b, otherwise $t_0 \leftarrow 0$

bne \$t0, \$zero, Less

go to Less if t0 != 0

Less:

3

1

Case/Switch Statement_{1/3}

•Switch statement allows the programmer to select one of many alternatives depending on a single value

•Can be implemented as a series of if -then-else statements (not efficient)

•Alternatively (or more efficiently)

First instruction within each case sequence is assigned a label (the label holds the address of the first instruction in the sequence)

>These addresses are stored in a jump table in memory

•Use new instruction

jr register (jump register)

Unconditional jump to the address specified in register

•Program loads appropriate entry from jump table into a register, then jump to proper address using a jump register

Case/Switch Statement_{3/3}

	slt \$t3, \$s5, \$zero	# test if $k < 0$
	bne \$t3, \$zero, Exit	# if k < 0 go to Exit
	slt \$t3,\$s5,\$t2	# test if k < 4
	beq \$t3, \$zero, Exit	# if $k \ge 4$ go to Exit
	add \$t1,\$s5,\$s5	# \$t1 ← 2 *k (Why?)
	add \$t1,\$t1,\$t1	# \$t1 ← 4 *k
	add \$t1,\$t1,\$t4	# \$t1 ← address of JumpTable[k]
	lw \$t0,0(\$t1)	# \$t0 ← JumpTable[k]
	jr \$tO	# jump to address found in register \$t0
L0:	add \$s0, \$s3, \$s4	# $(k = 0)$ then $f = i + j$
	j Exit	# break (go to Exit)
L1:	add \$\$0,\$\$1,\$\$2	#(k = 1) then $f = g+h$
	j Exit	# break (go to Exit)
L2:	sub \$s0,\$s1,\$s2	#(k=2) then $f = g-h$
	j Exit	
L3:	sub \$s0,\$s3,\$s4	#(k=3) then f = i- j
Exit:		# End of switch statement

Case/Switch Statement_{2/3}

switch (k) {

case 0: f = i + j; break; case 1: f = g + h; break; case 2: f = g - h; break; case 3: f = i - j; break;

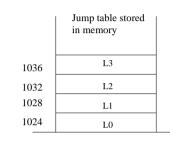
Variables f through k associated with \$s0 \rightarrow \$s5. Register \$t2 contains 4. Register \$t4 contains address of jump table in memory. MIPS Assembly code?

\$t4

1024

What we know so far

Address of instruction $f = i + j \rightarrow label L0$ Address of instruction $f = g + h \rightarrow label L1$ Address of instruction $f = g - h \rightarrow label L2$ Address of instruction $f = i - j \rightarrow label L3$



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ement_{3/3}

5

•Fig. 3.9 page 131
•instruction format for unconditional jump j (J-format)
<opcode (6 bits), target address (26 bits)> (more on target address in section 3.8)
•The jr instruction → (R-format)

opcode = 0, funct = 8, rs is jump register

•The slt instruction → (R-format)

Opcode = 0, funct = 42, rs, rd, rt

•The beq, bne instructions → (I-format)
•\$zero always contains the value 0
•The instruction beq \$s1, \$s2, 100 is an I-format instruction. The resulting machine code

•Op| rs | rt | address
4 | 17 | 18 | 25

•Why the value in address field is 25 and not 100? Will discuss this when we cover section 3.8

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